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050069

VO Netzwerktechnologie für Multimedia Anwendungen

Lecture 3: Review of Networking Concepts,
Multimedia Networking

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Endowed by

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Chapter 2: Revision of Networking Concepts

- Overview:
- Protocol layering and Internet protocol stack
- Circuit switching vs. packet switching
- Connectionless vs. connection-oriented networks, routing, forwarding, and switching
- Transport layer protocols
- **Application layer**
 - Sockets
 - Client-server and peer-to-peer communication
- Web services

Sockets 1/2

- **Process**
 - Program running within a host
 - Processes within same host communicate using **inter-process communication** (defined by OS).
 - Processes in different hosts communicate by exchanging **messages**
 - Application: process in a hosts
 - **Client process**: process that initiates communication
 - **Server process**: process that waits to be contacted
- Several processes running on the same host
 - **Identification** through 16 bit port numbers
 - Example port numbers:
 - HTTP server: 80
 - Mail server: 25

Sockets 2/2

- **Host**
 - Identification through unique 32 bit IP address
- **Socket**
 - End-point of an Internet Protocol-based communication
 - Components
 - Protocol (TCP, UDP, raw IP)
 - Local IP address
 - Local port
 - Remote IP address
 - Remote port
 - The remote address can be either
 - any valid IP address, or
 - 0.0.0.0 for listening sockets, or
 - 255.255.255.255 for broadcasting sockets
 - Process sends/receives messages to/from its socket

Application Layer Protocols 1/2

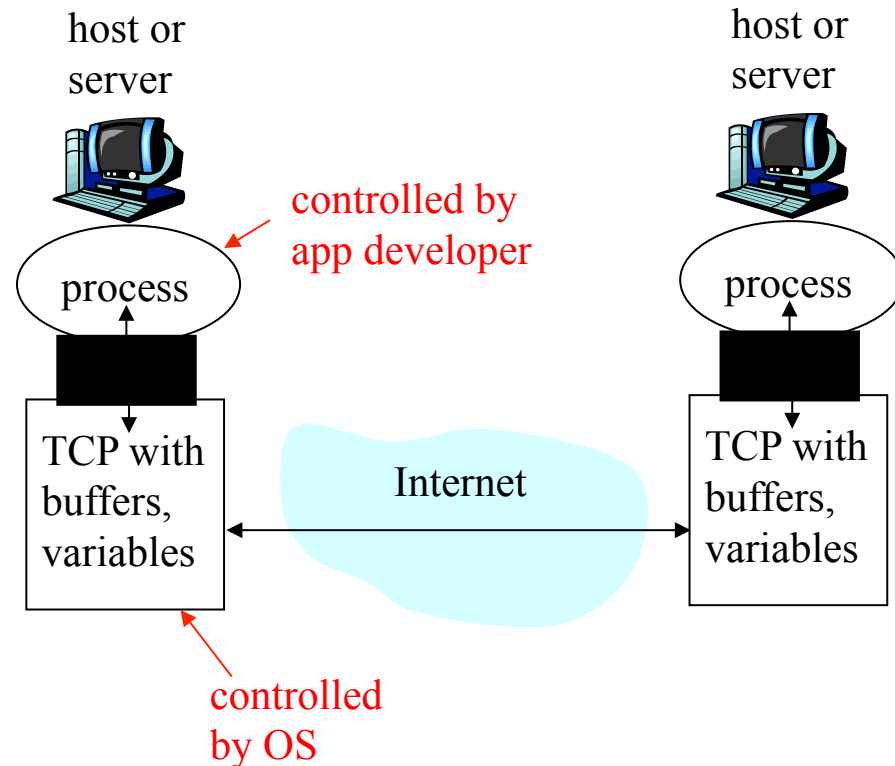
- Application layer protocols
 - Run on different end systems
 - Communicate over a network
 - e.g., Web: Web server software communicates with browser software
 - No interaction from core network devices required
 - New software only in end-points needed: rapid deployment possible
- Public-domain protocols
 - defined in RFCs
 - allows for interoperability
 - E.g., HTTP, SMTP
- Proprietary protocols
 - E.g., KaZaA

Application Layer Protocols 2/2

- Application layer protocols define
 - Types of messages exchanged
 - E.g., request & response messages
 - Syntax of message types
 - What fields in messages & how fields are delineated
 - Semantics of the fields
 - State machines
 - Rules for when and how processes send & respond to messages
 - Preferred transport protocol for msgs
 - Preferred port numbers to be contacted
 - E.g., http on port 80

Network Applications

- Create a new network application
 - Define new protocol
 - Deploy software and start it
 - Start communication by contacting those hosts on the right port
- Some network applications
 - E-mail
 - Web
 - Instant messaging
 - Remote login
 - P2P file sharing
 - Multi-user network games
 - Streaming stored video clips
 - Internet telephone
 - Real-time video conference
 - Massive parallel computing



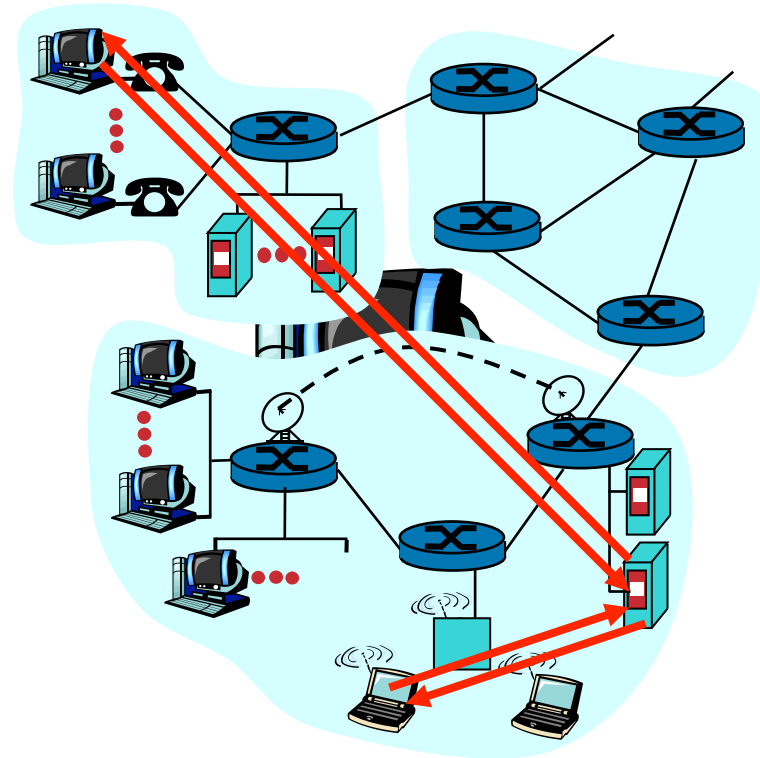
Client-Server Architecture

Server

- Always-on host
- Permanent IP address
- Server farms for scaling

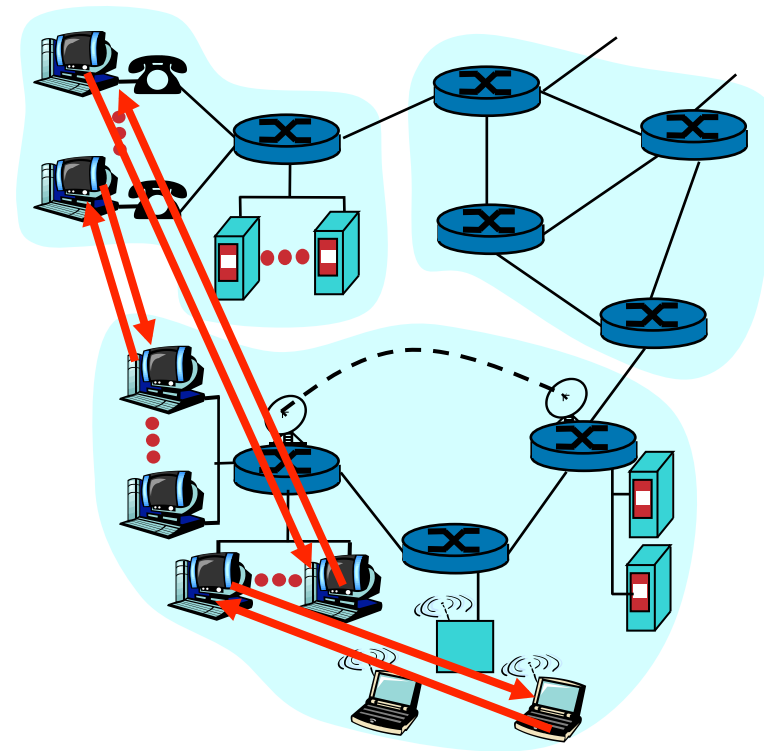
Clients

- Communicate with server
- May be intermittently connected
- May have dynamic IP addresses
- Do not communicate directly with each other



P2P Architecture

- Arbitrary end systems directly communicate
- Peers are intermittently connected and change IP addresses
- Note: applications with P2P architectures have client processes & server processes
- **Highly scalable**
- **But difficult to manage**



Was ist P2P?

- Def.: Ein Peer-to-Peer System ist ein selbstorganisierendes, verteiltes System aus miteinander verbundenen, gleichen und autonomen Knoten, zur gemeinsamen Nutzung von verteilten Ressourcen in einem Netzwerk ohne eine zentrale Instanz.
- Peer (engl. Gleichgestellter)
- Peer-to-Peer vs. Server-Client Prinzip
- Charakteristika von P2P Systemen:
 - Peers sind sowohl Clients als auch Server
 - Dezentralisierung
 - Nutzung nicht benötigter Ressourcen
 - Transient Connectivity
 - Autonomie der Peers
 - Keine Globale Sicht

Bewertung von P2P

- Vorteile:
 - Skalierbarkeit
 - Gemeinsame Nutzung von Ressourcen
 - Robuster gegen den zufälligen Ausfall einzelner Komponenten
- Probleme:
 - Peers und Verbindungen sind nicht zuverlässig
 - Hoher Kommunikationsaufwand

Anwendung von P2P

- Aktuelle Anwendungsbereiche
 - Filesharing (Bittorrent, KaZaa, Napster, eDonkey)
 - Instant-Messaging (Skype, CSpace)
 - Data Storage (PAST, OceanStore, FarSite)
 - Grid/Distributed Computing (Seti@Home, Folding@Home)
 - Collaboration / Groupware (Groove)
 - Kontrolle von Netzen (Tutschku et al., 2003 –)

P2P Overlay Network

- Overlay Network ist ein logisches Netz
- Setzt auf einem physikalische Netz (Internet) auf
- Besteht aus allen teilnehmenden Peers (Knoten)
- Jeder Knoten speichert eine Liste von Nachbarknoten
- Knoten sind durch (virtuelle) Kanten miteinander verbunden



P2P Klassifikation

Client-Server	Peer-to-Peer			
	Unstructured P2P			Structured P2P
	1. Generation		2. Generation	3. Generation
	Centralized P2P	Pure P2P	Hybrid P2P	DHT-based
<i>www</i>	<i>Napster</i>	<i>Gnutella 0.4, Freenet</i>	<i>Gnutella 0.6</i>	<i>Cord, CAN, Pastry</i>

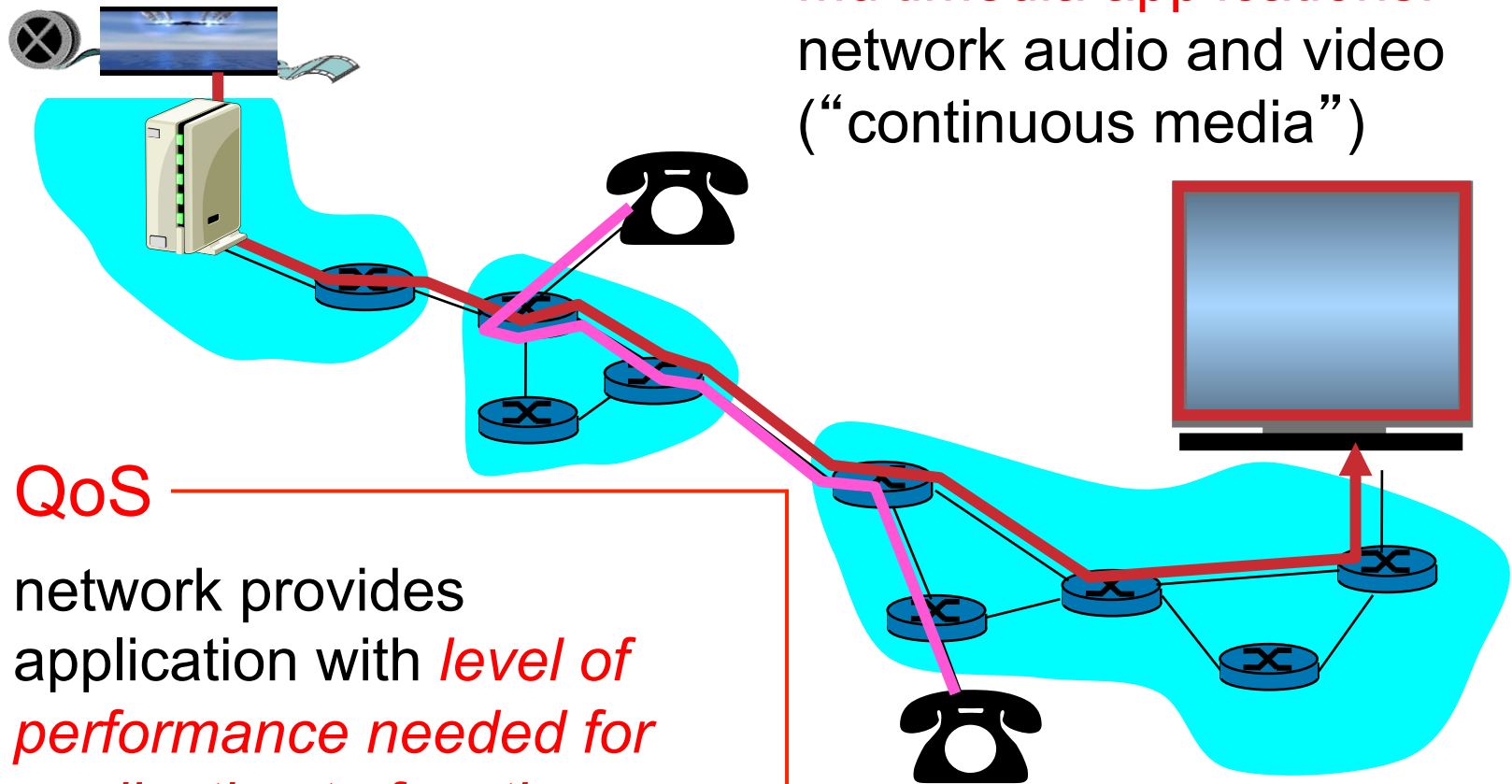
Chapter 3: Multimedia Networking

Overview

- ▶ **3.1 Multimedia Networking Applications**
- ▶ 3.2 Streaming stored audio and video
- ▶ 3.3 Real-time Multimedia: Internet Phone study
- ▶ 3.4 Protocols for Real-Time Interactive Applications
 - RTP, RTCP
- ▶ 3.5 IP Telephony, SIP, and H.323
- ▶ 3.6 Distributing Multimedia: content distribution networks

Multimedia, Quality of Service: What is it?

Multimedia applications:
network audio and video
("continuous media")



QoS

network provides
application with *level of
performance needed for
application to function.*

MM Networking Applications

Classes of MM applications:

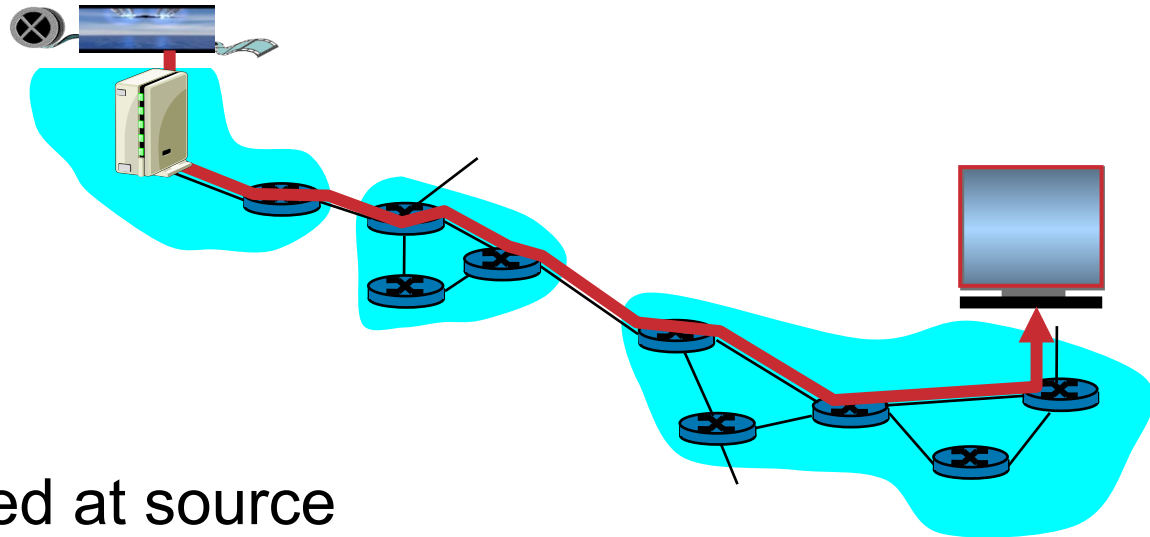
- Streaming stored audio and video
- Streaming live audio and video
- Real-time interactive audio and video

Fundamental characteristics:

- Typically **delay sensitive**
 - end-to-end delay
 - delay jitter
- But **loss tolerant**: infrequent losses cause minor glitches
- Antithesis of data, which are loss intolerant but delay tolerant

Jitter is the variability of packet delays within the same packet stream

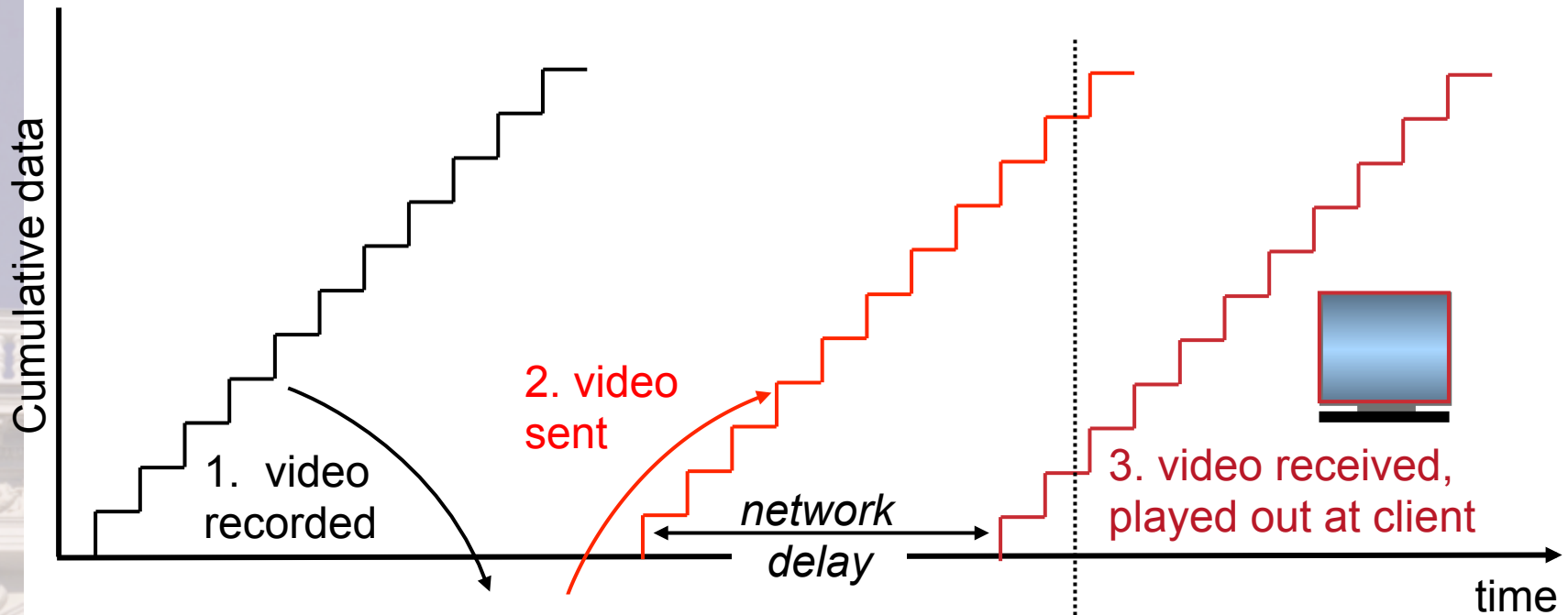
Streaming Stored Multimedia



Streaming:

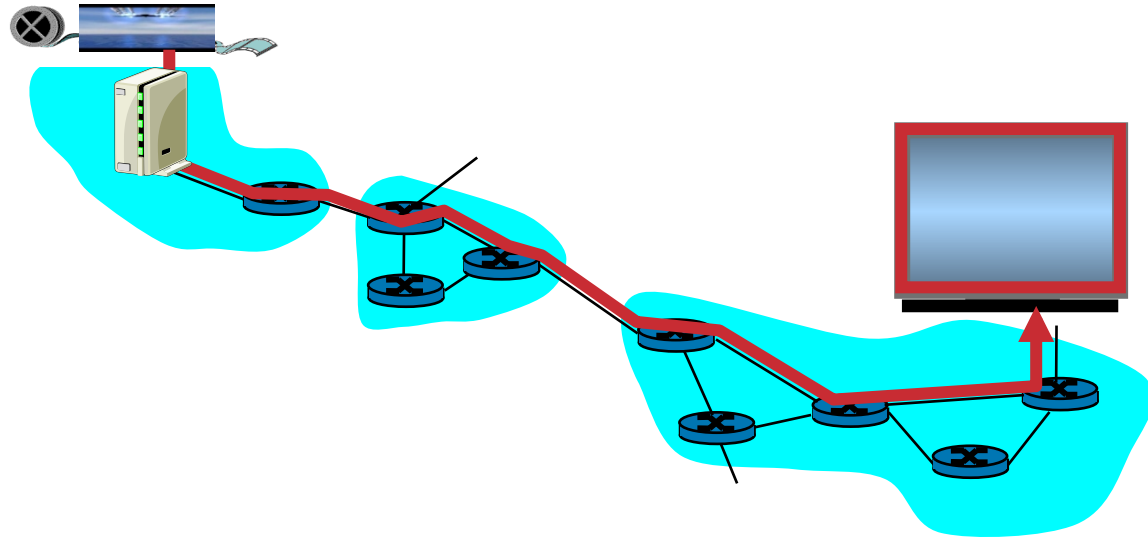
- media stored at source
- transmitted to client
- streaming: client playout begins *before* all data has arrived
- timing constraint for still-to-be transmitted data: in time for playout

Streaming Stored Multimedia: What is it?



streaming: at this time, client playing out early part of video, while server still sending later part of video

Streaming Stored Multimedia: Interactivity



- *VCR-like functionality*: client can pause, rewind, FF, push slider bar
 - 10 sec initial delay OK
 - 1-2 sec until command effect OK
 - RTSP often used (more later)
- timing constraint for still-to-be transmitted data: in time for playout

Examples:

- Internet radio talk show
- Live sporting event

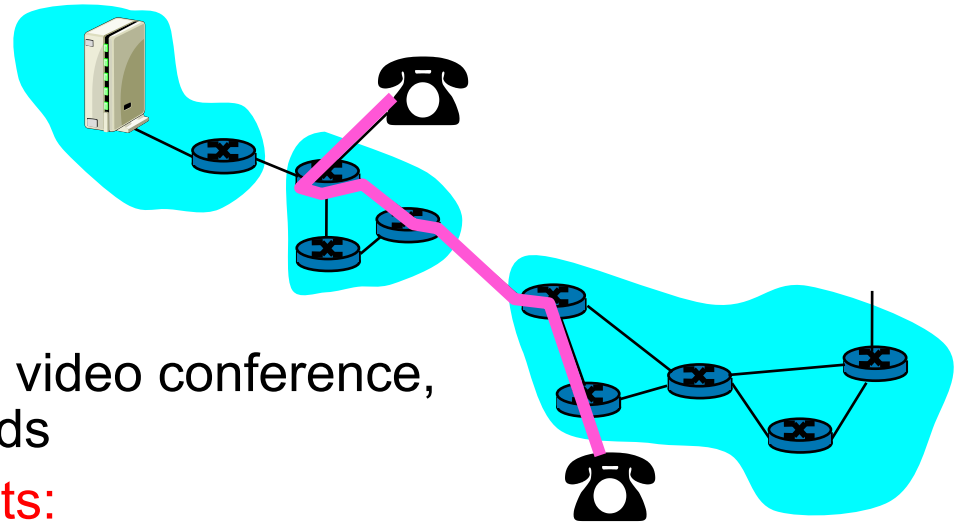
Streaming

- playback buffer
- playback can lag tens of seconds after transmission
- still have timing constraint

Interactivity

- fast forward impossible
- rewind, pause possible!

Interactive, Real-Time Multimedia



- **applications:** IP telephony, video conference, distributed interactive worlds
- **end-end delay requirements:**
 - audio: < 150 msec good, < 400 msec OK
 - includes application-level (packetization) and network delays
 - higher delays noticeable, impair interactivity
- **session initialization**
 - how does callee advertise its IP address, port number, encoding algorithms?

Multimedia Over Today's Internet

TCP/UDP/IP: “best-effort service”

- *no* guarantees on delay, loss



But you said multimedia apps requires QoS and level of performance to be effective!



Today's Internet multimedia applications use application-level techniques to mitigate (as best possible) effects of delay, loss

How should the Internet evolve to better support multimedia?

Integrated services philosophy:

- Fundamental changes in Internet so that apps can reserve end-to-end bandwidth
- Requires new, complex software in hosts & routers

Laissez-faire

- no major changes
- more bandwidth when needed
- content distribution, application-layer multicast
 - application layer

Differentiated services philosophy:

- Fewer changes to Internet infrastructure, yet provide 1st and 2nd class service.



What's your opinion?

A few words about audio compression

- Analog signal sampled at constant rate
 - telephone: 8,000 samples/sec
 - CD music: 44,100 samples/sec
- Each sample quantized, i.e., rounded
 - e.g., $2^8=256$ possible quantized values
- Each quantized value represented by bits
 - 8 bits for 256 values
- Example: 8,000 samples/sec, 256 quantized values --> 64,000 bps
- Receiver converts it back to analog signal:
 - some quality reduction

Example rates

- CD: 1.411 Mbps
- MP3: 96, 128, 160 kbps
- Internet telephony: 5.3 - 13 kbps

A few words about video compression

- Video is sequence of images displayed at constant rate
 - e.g. 24 images/sec
 - Digital image is array of pixels
 - Each pixel represented by bits
 - Redundancy
 - spatial
 - temporal
- Examples:**
- MPEG 1 (CD-ROM) 1.5 Mbps
 - MPEG2 (DVD) 3-6 Mbps
 - MPEG4 (often used in Internet, < 1 Mbps)
- Research:**
- Layered (scalable) video
 - adapt layers to available bandwidth



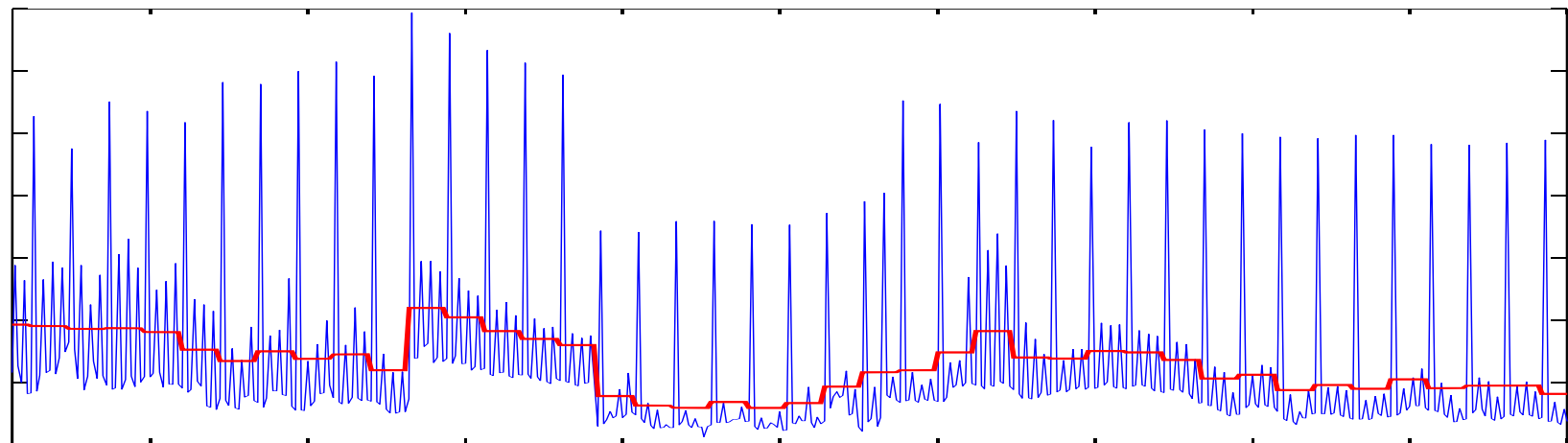
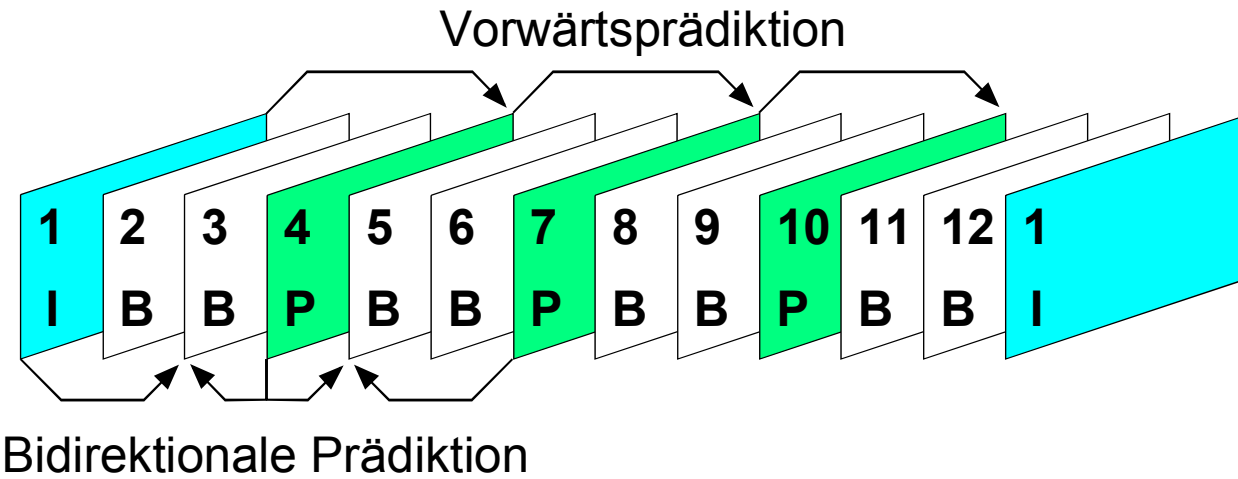
Generation of MPEG Video Sequences

- Given: sequence of digital images
- MPEG compression is combination of
 - Intra-frame compression (spatial redundancy reduction)
 - Discrete-Cosine Transformation (DCT):
8x8 pixel blocks \Rightarrow DCT \Rightarrow 8x8 DC coefficients
 - Quantization
 - Zig-zag entropy encoding
 - Inter-frame compression (temporal redundancy reduction)
 - Block-based motion compensation

Inter-Frame Dependencies

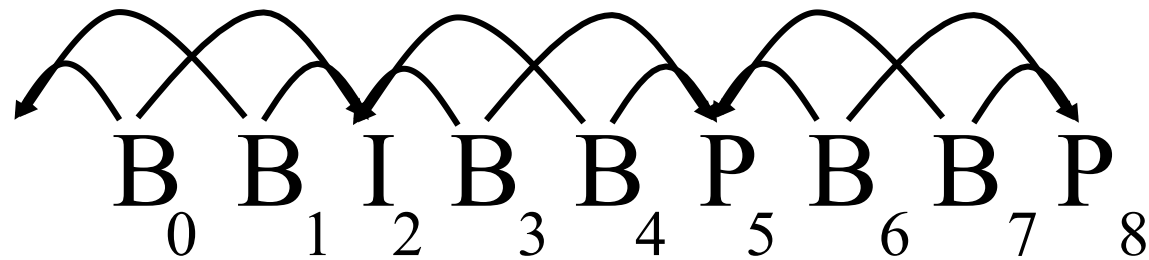
- Frame types
 - Intracoded frames (I-frame)
 - Does not depend on any other frames
 - Most important information, largest frame size
 - Predicted frames (P-frame)
 - Depends on preceding I- or P-frame
 - Medium frame size
 - Bidirectional frames (B-frame)
 - Depends on preceding and succeeding I- or P-frame
 - Small frame size, no other frame depends on it
- Group of Pictures (GoP)
 - All frames following and depending on a specific I-frame

Measured Time Series of MPEG Frame Sizes



Display and Bitstream Order

- Display order



- Bitstream order

